

Upper bridge & Lower bridge start in middle.

Data structure for CH.

Circular Doubly-linked list.

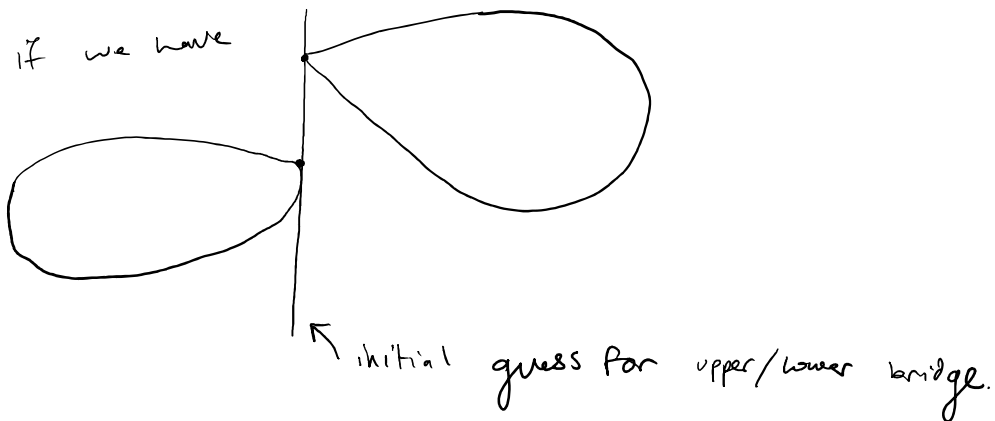
Suppose we have 3 pts:

$$P_i = (x_i, y_i)$$

(P_1, P_2, P_3) in CCW order iff

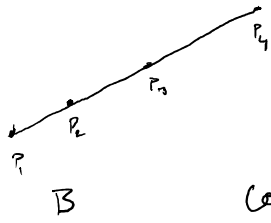
$$\begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix} > 0$$

CW if < 0 , Collinear if $= 0$.



just add a case for if line is vertical.

or avoid dividing along this line.



convex hull is just $\overline{P_1 P_4}$

Dynamic Programming

Optimization problem

: solution satisfies condition & minimizes objective function.
or maximizes

Problems & subproblems:

$P(i, j)$ denotes problem of finding closest pair in $A_{i,j} = \{P_i, \dots, P_j\}$
where $1 \leq i \leq j \leq n$

original problem is $P(1, n)$

$A = \left\{ P(i, j) : 1 \leq i \leq j \leq n \right\}$ is one class of problems.

we could also let $P(i)$ denote finding closest pair in $\{P_i, \dots, P_i\}$

$B = \{P(i) : 1 \leq i \leq n\}$

$$B \subset A$$

Problem: construct optimal solution (x_1, \dots, x_n)

options for x_i : op_1, op_2, \dots, op_d

each option leads to a subproblem P_j :

given $x_i = op_j$, find optimal solution $(op_j, x_{2j}, \dots, x_{nj})$

The best of these j optimal sub solns is optimal soln

DP only works if P_j is a problem similar to the original problem.

Now have an exponential tree of problems.

if many problems are the same and have few parameters, exponential tree is reduced to polynomial.

w/ DAC subproblems are disjoint.