

- Quiz:
- 1 True ✓
  - 2  $\Theta(2^{4n^2} + 3n)$  ✓
  - 3  $\Theta(n^5)$

## Math foundations

Analysis: runtime := count of basic operations.

$T(n)$  := worst case runtime.

↳ sometimes write  $T_w(n)$

for sequential search,  $T(n) = 3n + 1$  or something  $\in O(n)$  ↖ or  $\Theta$  or  $\Omega$

$$O(g) = \left\{ f \in 2^{\mathbb{N}} \mid \exists c > 0, n_0 \text{ s.t. } f(n) \leq cg(n) \quad \forall n \geq n_0 \right\}$$

$$\Omega(g) = \left\{ f \in 2^{\mathbb{N}} \mid \exists c > 0, n_0 \text{ s.t. } f(n) \geq cg(n) \quad \forall n \geq n_0 \right\}$$

$$\Theta(g) = O(g) \cap \Omega(g)$$

Theorem: If  $f_1 \in O(g_1)$  and  $f_2 \in O(g_2)$  then

$$f_1 + f_2 \in O(g_1 + g_2) = O(\max(g_1, g_2))$$

$$f_1 \cdot f_2 \in O(g_1 \cdot g_2)$$

Proof: easy constant building.

same theorem for  $\Omega$ .

Perhaps better definitions:  $\rightarrow$  Using more words and fewer constants

$$O(g) = \left\{ f \in \mathbb{Z}^{\mathbb{N}} \mid \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \neq \infty \right\}$$

$$\Omega(g) = \left\{ \text{"} \quad \quad \quad \text{"} \neq 0 \right\}$$

$$\Theta(g) = O(g) \cap \Omega(g)$$

$$o(g) = \left\{ f \in \mathbb{Z}^{\mathbb{N}} \mid \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \right\}$$

$$\omega(g) = \left\{ \text{"} \quad \quad \quad \text{"} = \infty \right\}$$

$$\theta(g) = \emptyset$$

If  $f \in O(g)$  and  $g \in O(h)$  then  $f \in O(h)$

same w/  $\Theta, \Omega$ .

for  $g: \mathbb{N}^k \rightarrow \mathbb{N}$ ,  $O(g) = \left\{ f: \mathbb{N}^k \rightarrow \mathbb{N} \mid \exists c, r > 0 \text{ s.t. } f(\vec{n}) \leq c g(\vec{n}) \text{ when } |\vec{n}| > r \right\}$

$$4m^2 + 9mn^2 + 6n^3 \notin O(m^2 + n^3)$$

since when  $m = n^2$ ,  $mn^2 \in \omega(m^2 + n^3)$

2)  $\sqrt{n}$