

$$R_{e,B} : \chi_A \neq \Phi_e^B$$

$$R_{e,C} : \chi_A \neq \Phi_e^C$$

Lemma 2: let  $m = \lim_s L(1, B, s)$ . Then at least one  $y \leq m$  is a permanent witness for  $R_{1,B}$ , i.e.  $A(y) \neq \Phi_y^B$ .

PF let  $s$  be the first stage at which  $m$  is hit.

If  $\exists y < m$  that enters  $A$  after stage  $s$ ,

then  $y$  is a permanent witness since  $\forall t \geq s, [\Phi_y^B(t)] \downarrow = 0$ .

So assume that no  $y < m$  enters  $A$  after stage  $s$ .

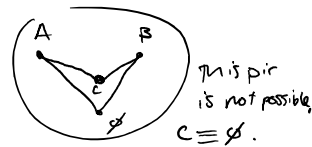
Then  $m$  is the permanent witness since  $m$  is the high-water mark,

so if it were not a permanent witness we'd get a contradiction.  $\square$

We argue by induction on the priority that each requirement acts only finitely often.

New theorem

DFN: A pair of non-recursive sets  $A$  and  $B$  is minimal if  $(\forall C) [(C \leq_T A \ \& \ C \leq_T B) \implies C \text{ is recursive}]$ .



Lachlan-Yates Thm: There is a minimal pair of r.e. sets.

Proof: We'll construct  $A$  and  $B$  in stages to satisfy the following requirements:

$$R_e : \bar{A} \neq W_e, \quad Q_e : \bar{B} \neq W_e, \quad N_e : (\Phi_e^A = \Phi_e^B = \text{some true fn } f) \implies f \text{ is computable.}$$